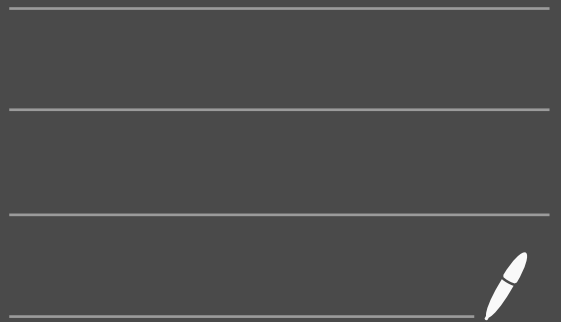


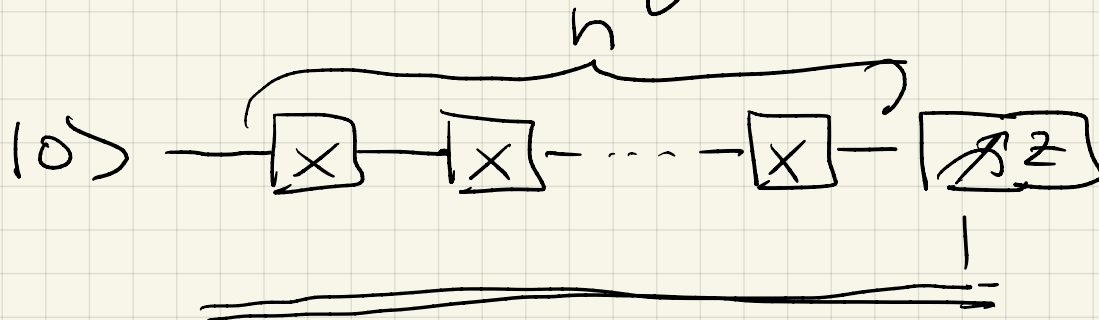
# Lecture 2

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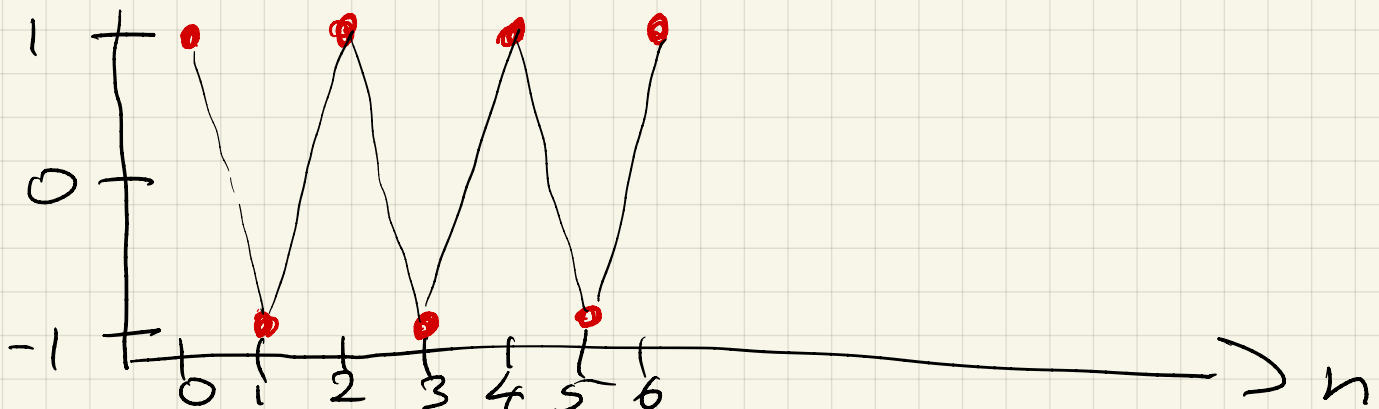


# Noisy Simulations

- From yesterday's lecture saw that noise can severely affect the results we can obtain from quantum computers.
- Goal of today's lecture is to understand the origin of this noise better and find ways to mitigate this noise.
- To start, consider a very simple circuit



- What do we expect from this circuit?



- In practice, no quantum computer is perfect, and there are 3 kinds of errors that a computer can make:

1) Coherent error:

The operations performed are not exact, so instead of applying  $X$ -gate, applies some slightly different  $X$  gate

2) Sampling error:

Quantum mechanics is probabilistic, and one therefore has statistical uncertainties

3) Measurement error:

The probability to measure a pure state  $|\psi\rangle = |0,1\rangle$  in  $|0,1\rangle$  state is not 100%

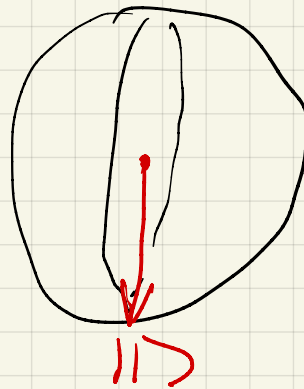
4) Incoherent error:

The quantum computer is coupled to environment, which will lead to decoherence.

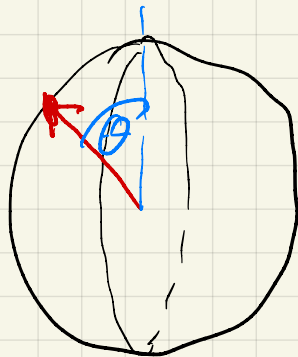
• Show simple circuit in Qiskit.

# Coherent errors

- How does a QC implement  $X$ -gate?
- Know that any 1-qubit state lies on Bloch-sphere



- Can move states on Bloch-sphere by applying some pulses to my qubits
- One possibility is  $X$ -rotation  $R_x(\theta)$



- $\hat{X}|0\rangle = |1\rangle$  ,  $\hat{X}|1\rangle = |0\rangle$
- This is the same as  $R_x(\pi)$



- To be precise  $\hat{X} = i \hat{R}_x(\pi)$
- So when applying an  $\hat{X}$  gate, tell QC "Apply  $R_x$  rotation with angle  $\pi$ "
- But can never do things precisely, so will typically do  $\hat{R}_x(\pi + \epsilon) \equiv \hat{X}$
- Rotating by  $\pi + \epsilon$  is same as first rotating by  $\pi$  and then rotating by  $\epsilon$

$$\Rightarrow \hat{X} = R_x(\epsilon) \hat{X}$$

- Using simple math

$$R_x(\epsilon) = \exp\left(-\frac{i\epsilon}{2} \hat{X}\right)$$

$$= 1 - \frac{i\epsilon}{2} \hat{X} + \frac{1}{2!} \left(\frac{i\epsilon}{2}\right)^2 \hat{X}^2 + \frac{1}{3!} \left(\frac{i\epsilon}{2}\right)^3 \hat{X}^3 + \dots$$

$$= 1 - \frac{i\epsilon}{2} \hat{X} + \frac{1}{2!} \left(\frac{i\epsilon}{2}\right)^2 + \frac{1}{3!} \left(\frac{i\epsilon}{2}\right)^3 \hat{X} + \dots$$

$$= \cos\left(\frac{\epsilon}{2}\right) I - i \sin\left(\frac{\epsilon}{2}\right) \hat{X}$$

$$\Rightarrow \hat{X} = \cos\left(\frac{\epsilon}{2}\right) \hat{X} - i \sin\left(\frac{\epsilon}{2}\right) I$$

In general

$$\hat{X}^n = R_x^n(\epsilon) \hat{X}^n$$

$$= R_x(n\varepsilon) \hat{X}^n$$

$$= \cos\left(\frac{n\varepsilon}{2}\right) \hat{X}^n - i \sin\left(\frac{n\varepsilon}{2}\right) \hat{X}^{n+1}$$

- Show results with different values of  $\varepsilon$ .

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## Sampling error

- Know that quantum mechanics is probabilistic, so only recover observables probabilistically.
- Say  $|\psi\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$   
 $\Rightarrow |\langle 0|\psi\rangle|^2 = \frac{1}{2}$  ,  $|\langle 1|\psi\rangle|^2 = \frac{1}{2}$
- But if only doing a single measurement will find either  $|0\rangle$  or  $|1\rangle$  when measuring.
- Probabilities only recovered statistically
- All know that statistical uncertainty scales as  $1/\sqrt{N}$

◦ In general

$$|\psi\rangle = \alpha |0\rangle + \sqrt{1-\alpha^2} |1\rangle$$

$$\Rightarrow \sigma = \sqrt{\frac{\alpha^2(1-\alpha^2)}{N}}$$

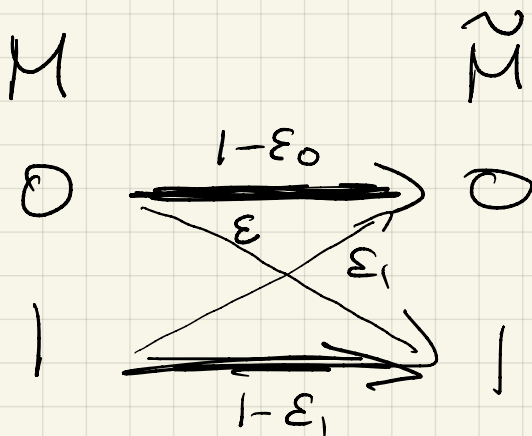
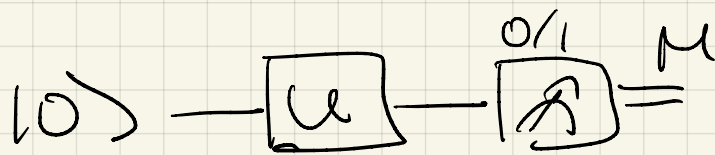
$\Rightarrow$  need large number of measurements

◦ error maximal if  $\alpha = \frac{1}{\sqrt{2}}$ .

◦ Show with more or less slots

## Measurement error

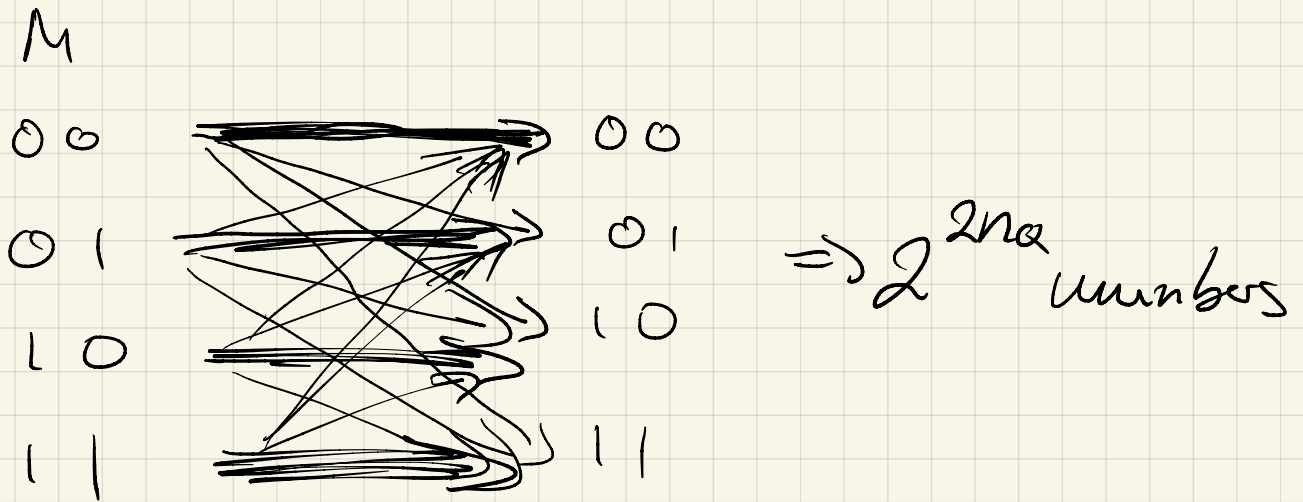
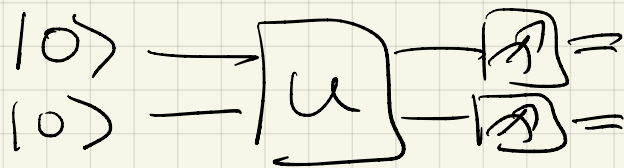
◦ Assume we have circuit



2 numbers describe measurement error of each qubit.

$\Rightarrow 2n_Q$  numbers

- In general, can have measurement errors that depend on all qubits



- Show results for different measurement error values

## Incoherent error

- Incoherent errors arise from interactions with environment.
- The combined system of  $|q\rangle$  and  $|e\rangle$  undergoes unitary evolution

- A general way of writing unitary evolution of  $|\psi\rangle \otimes |e\rangle$  can be written as

$$|\psi\rangle \otimes |e\rangle \xrightarrow{U} \sum_k [E_k |\psi\rangle] \otimes |e_k\rangle, \quad \sum_k E_k^\dagger E_k = 1$$

- Density matrix transforms as

$$\rho = |\psi\rangle\langle\psi| \otimes |e\rangle\langle e| \xrightarrow{U} \sum_{k,l} [E_k |\psi\rangle\langle\psi| E_l^\dagger] \otimes |e_k\rangle\langle e_l|$$

- After tracing out environment

$$\begin{aligned} \rho_\psi &\rightarrow \rho_\psi = \text{Tr}_e[\rho] \\ &= \sum_{k,l} [E_k |\psi\rangle\langle\psi| E_l^\dagger] \langle e_l | e_k \rangle \\ &= \sum_k [E_k |\psi\rangle\langle\psi| E_k^\dagger] \end{aligned}$$

- Note that from derivation above,  $k$  runs over all basis state in  $|e\rangle$ .
- But we know that operators  $E_k$  act on state  $|\psi\rangle$  so are  $d \times d$  matrices.  $\Rightarrow$  number restricted by dimension of  $|\psi\rangle$

$$\begin{aligned} \bullet \quad E_k &= \sum_i c_{ik} B_i, \quad \text{Tr}[B_i B_j] = d \delta_{ij} \\ \Rightarrow \sum_k E_k \rho E_k^\dagger &= \sum_{ij} \left( \sum_k c_{ik} c_{jk}^* \right) B_i \rho B_j^\dagger \end{aligned}$$

$$= \sum_{ij} d_{ij} B_i \rho B_j^\dagger$$

$$d_{ij} = \sum_m U_{im} P_m U_{jm}^* , \quad A_m = \sum_i U_{im} B_i$$

$$= \sum_{m=1}^{d^2} P_m A_m \rho A_m^\dagger$$

- From original unitarity,  $\sum_m P_m = 1$

$d$ : dim of Hilbert space

- All interactions with environment can be parametrized by  $d^2$  Krauss operators.

- Consider a single qubit  $\Rightarrow d=2$

$\Rightarrow$  4 Krauss operators

- Choose as  $\mathbb{1}, X, Y, Z$

$$\rho \rightarrow P_0 \rho + P_x X \rho X + P_y Y \rho Y + P_z Z \rho Z$$

bit-flip                  bitphase-flip                  phase-flip

- Depolarizing channel assumes

$$P_x = P_y = P_z = \frac{P}{4}$$

$$\Rightarrow \rho \rightarrow (1 - \frac{3}{4}P) \rho + \frac{P}{4} \cdot X\rho X + \frac{P}{4} Y\rho Y + \frac{P}{4} Z\rho Z$$

$$= (1 - P) \rho + \frac{P}{4} (\mathbb{1}\rho\mathbb{1} + X\rho X + Y\rho Y + Z\rho Z)$$

Take a generic density matrix  $\rho = \begin{pmatrix} a & b \\ c & 1-a \end{pmatrix}$

$$\mathbb{1}\rho\mathbb{1} = \begin{pmatrix} a & b \\ c & 1-a \end{pmatrix}, \quad X\rho X = \begin{pmatrix} 1-a & c \\ b & a \end{pmatrix}$$

$$Y\rho Y = \begin{pmatrix} 1-a & -c \\ -b & a \end{pmatrix}, \quad Z\rho Z = \begin{pmatrix} a & -b \\ -c & 1-a \end{pmatrix}$$

$$\Rightarrow \mathbb{1}\rho\mathbb{1} + X\rho X + Y\rho Y + Z\rho Z = \begin{pmatrix} 2 & 0 \\ 0 & 2 \end{pmatrix}$$

$$\Rightarrow \rho \rightarrow (1 - P) \rho + \frac{P}{2} \mathbb{1}$$

• Show result with depolarizing channel

• Show results with all channels.

# Mitigating noise

- How big are typical noise levels?
- Assume incoherent errors for now
- Noise classified through probabilities in Kraus decomposition.
- Noise on single qubits  $\sim 10^{-4} - 10^{-5}$
- Noise on entangling (CNOT) gates  $\sim 10^{-2} - 10^{-3}$   
 $\Rightarrow$  focus on CNOT gates
- Combining unitary operation with depolarizing channel gives

$$\rho \rightarrow (1 - \epsilon) U_q \rho U_q + \epsilon \left( \frac{I_q}{2^{n_q}} \otimes \rho_q \right)$$

- Apply this to CNOT gate acting on qubits  $k$  and  $l$   $U_c^{(kl)}$

$$\rho \rightarrow (1 - \epsilon^{(kl)}) U_c^{(kl)} \rho U_c^{(kl)} + \epsilon^{(kl)} \left( \frac{I_4}{4} \otimes \rho_{kl} \right)$$



• Apply 2 CNOT operations

$$\rho \xrightarrow{2x} (1 - \varepsilon^{(ke)}) U_C \left[ (1 - \varepsilon^{(ke)}) U_C \rho U_C^{(ke)} + \varepsilon^{(ke)} \left( \frac{I_4}{4} \oplus \rho_{ke} \right) \right]$$

$$+ \varepsilon^{ke} \left( \frac{I_4}{4} \oplus \rho_{ke} \right)$$

$$= (1 - \varepsilon^{(ke)})^2 U_C^{(ke)} U_C^{(ke)} \rho U_C^{(ke)} U_C^{(ke)}$$

$$+ (1 - \varepsilon^{(ke)}) \varepsilon^{(ke)} \frac{U_C^{(ke)} \rho U_C^{(ke)}}{4} \oplus \rho_{ke}$$

$$+ \varepsilon^{(ke)} \left( \frac{I_4}{4} \oplus \rho_{ke} \right)$$

$$= (1 - \varepsilon^{(ke)})^2 \rho + \left[ \varepsilon^{(ke)} - \varepsilon^{(ke)2} + \varepsilon^{(ke)} \right] \left( \frac{I_4}{4} \oplus \rho_{ke} \right)$$

$$= (1 - \varepsilon^{(ke)})^2 \rho + \left[ 1 - (1 - \varepsilon^{(ke)})^2 \right] \left( \frac{I_4}{4} \oplus \rho_{ke} \right)$$

• Apply 3 CNOTS

$$\rho \xrightarrow{3x} (1 - \varepsilon^{(ke)})^3 U_C \rho U_C + \left[ 1 - (1 - \varepsilon^{(ke)})^3 \right] \left( \frac{I_4}{4} \oplus \rho_{ke} \right)$$

• Apply  $r = 2n + 1$  CNOTS

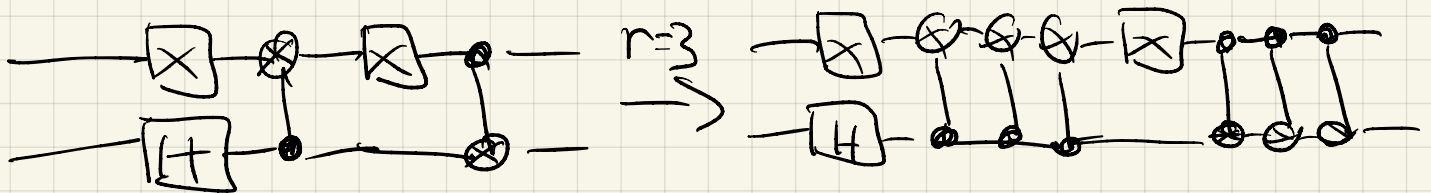
$$\rho \xrightarrow{rx} (1 - \varepsilon_i)^r U_C \rho U_C + \left[ 1 - (1 - \varepsilon)^r \right] \left( \frac{I_4}{4} \oplus \rho_{ke} \right)$$

- Taylor expanding around  $\epsilon = 0$

$$\rho \xrightarrow{\Gamma} (1 - \Gamma \sum_i \epsilon_i) U_c^{(ke)} \rho U_c^{(ke)} + \Gamma \sum_i \epsilon_i \left( \frac{+ \epsilon_i}{4} \oplus \rho_{ke} \right)$$

- Applying  $\Gamma_i$  CNOT gates is same as applying single CNOT, but with dephasing parameter amplified by  $\Gamma_i$

- Now consider a circuit parametrized by  $\Gamma$  (how many CNOTS for each CNOT)



- Circuit will produce density matrix  $\rho(\Gamma)$

- Now measure observable  $M$  and obtain

$$\langle M(\Gamma) \rangle = \text{Tr} [M \rho(\Gamma)]$$

- Use expression of  $\rho(\Gamma)$  to obtain

$$\langle M(\Gamma) \rangle = \left( 1 - \Gamma \sum_{i=1}^{N_c} \epsilon_i \right) \langle M \rangle_{ex} + \Gamma \sum_{i=1}^{N_c} \epsilon_i \langle M \rangle_{dep}$$

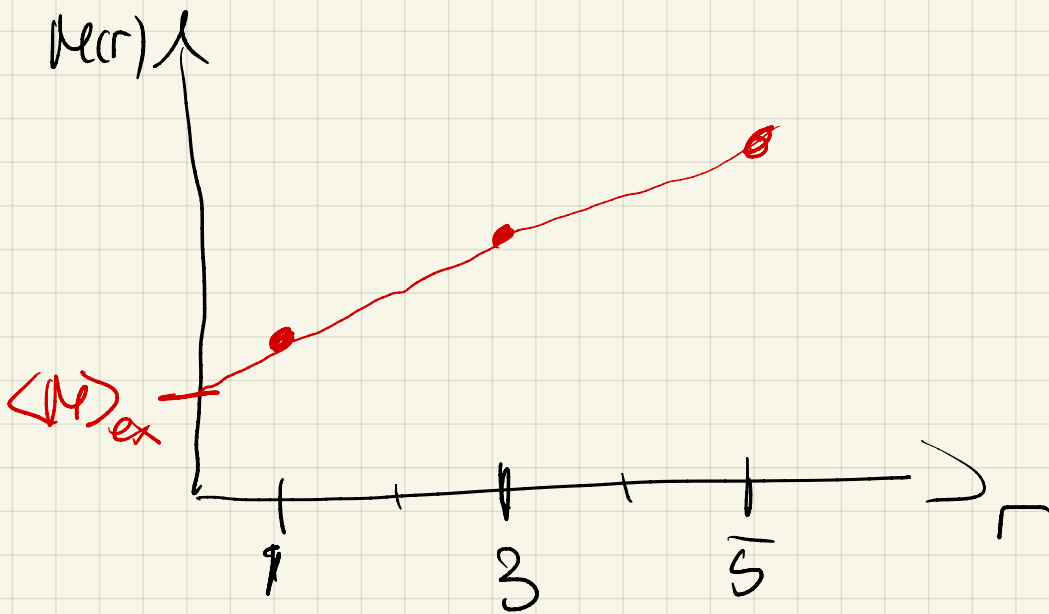
where  $\langle M \rangle_{ex}$  is expectation value in

absence of noise, and  $\langle M \rangle_{\text{dep}}$  is exp value  
if WOT; is replaced by depolarizing channel

• Thus

$$\langle M \rangle_{\text{ex}} = \langle M(\infty) \rangle$$

and one can extract  $\langle M \rangle_{\text{ex}}$  by extrapolating  
 $\langle M(r) \rangle$  to zero



• So far have done linear fits, which eliminated  
 $O(\epsilon^2)$  errors

• In general can eliminate depolarizing error  
to all orders

- Review linear fit again

$$\langle M \rangle(1) = \langle M \rangle_{\text{ex}} + N_{\text{out}} \varepsilon \left[ \sum_i \langle M \rangle_{\text{dep}_i} - \langle M \rangle_{\text{ex}} \right] + O(\varepsilon^2)$$

$$\langle M \rangle(3) = \langle M \rangle_{\text{ex}} + 3N_{\text{out}} \varepsilon \left[ \sum_i \langle M \rangle_{\text{dep}_i} - \langle M \rangle_{\text{ex}} \right] + O(\varepsilon^2)$$

$$\Rightarrow \frac{3}{2} \langle M \rangle(1) - \frac{1}{2} \langle M \rangle(3) = \langle M \rangle_{\text{ex}} + O(\varepsilon^2)$$

- This can be generalised (Richardson extrapolation)

- Assume we have  $\langle M \rangle(r)$  for  $r=1, 3, \dots, r_{\text{max}}$

- Write linear combination

$$\sum_{n=0}^{N_{\text{max}}} a_n \langle M \rangle(1+2^n) \stackrel{!}{=} \langle M \rangle_{\text{ex}} + O(\varepsilon^{N_{\text{max}}+1})$$

- One can work out general equation

$$\langle M \rangle(r) = (1-\varepsilon)^{N_c r} \langle M \rangle_{\text{ex}}$$

$$+ (1-\varepsilon)^{(N_c-1)r} \left[ 1 - (1-\varepsilon)^r \right] \sum_i \langle M \rangle_{\text{dep}_i}$$

$$+ (1-\varepsilon)^{(N_c-2)r} \left[ 1 - (1-\varepsilon)^r \right]^2 \sum_{i_1, i_2} \langle M \rangle_{\text{dep}_{i_1, i_2}}$$

+ ...

$$+ \left[ 1 - (1-\varepsilon)^r \right]^{N_c} \sum_{i_1, \dots, i_{N_c}} \langle M \rangle_{\text{dep}_{i_1, \dots, i_{N_c}}}$$

$$\begin{aligned}
&= \langle M \rangle_{ex} - f_{n_c}(r, \varepsilon) \langle M \rangle_{ex} \\
&\quad + [f_{n_c}(r, \varepsilon) - f_{n_c-1}(r, \varepsilon)] \sum_i \langle M \rangle_{dep_i} \\
&\quad + [f_{n_c}(r, \varepsilon) - f_{n_c-2}(r, \varepsilon)] \sum_{i_1, i_2} \langle M \rangle_{dep_{i_1, i_2}} \\
&\quad + \dots \\
&\quad + f_{n_c}(r, \varepsilon)^{n_c} \sum_{i_1, \dots, i_{n_c}} \langle M \rangle_{dep_{i_1, \dots, i_{n_c}}}
\end{aligned}$$

with  $f_n(r, \varepsilon) = 1 - (1 - \varepsilon)^{nr}$

- Important to remember that  $\langle M \rangle_{ex}$   $\langle M \rangle_{dep_{i_1}}$   $\langle M \rangle_{dep_{i_1, \dots, i_{n_c}}}$  are results of observables in noiseless circuit one does not have access to.
- But one can find coefficients  $a(i)$  such that all  $\langle M \rangle_{dep_i}$  values cancel to  $O(\varepsilon^{n_{max}+1})$
- One finds

$$a(i) = \prod_{j=0, j \neq i}^{n_{max}} \frac{(1+2j)}{2(j-i)}$$

$$= \frac{2^{-2n_{\max}}}{i!} \frac{(-1)^i}{1+2i} \frac{(1+2n_{\max})!}{n_{\max}! (n_{\max}-i)!}$$

◦ Note that this is the same as performing a polynomial fit with degree  $n_{\max} - 1$